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# Data Structures and Algorithms in Java™

Sixth Edition

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## Study Guide: Hints to Exercises

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# Chapter 10 Maps, Hash Tables, and Skip Lists

## Hints

### Reinforcement

- R-10.1)** The first insertion is  $O(1)$ , the second is  $O(2)$ , ...
- R-10.2)** Use the `PositionalList.remove` method to delete an entry from the map.
- R-10.3)** Take advantage of the existing `findIndex` method.
- R-10.4)** Think about which of the schemes use the array supporting the hash table exclusively and which of the schemes use additional storage external to the hash table.
- R-10.5)** There is a lot of symmetry and repetition in this string, so avoid a hash code that would not deal with this.
- R-10.6)** Try to mimic the figure in the book.
- R-10.7)** Try to mimic the figure in the book.
- R-10.8)** The failure occurs because no empty slot is found. For the drawing, try to mimic the figure in the book.
- R-10.9)** Try to mimic the figure in the book.
- R-10.10)** Think of the worst-case time for inserting every entry in the same cell in the hash table.
- R-10.11)** Mimic the way the figure is drawn.
- R-10.12)** Combine the hash values of the two components to make a hash value for the pair.
- R-10.13)** There is a subtle flaw—but can you find it?
- R-10.14)** The load factor can be controlled from within the abstract class, but there must be means for setting the parameter (either through the constructor, or a new method).
- R-10.15)** It is okay to insert a new entry on “top” of the deactivated entry object.

- R-10.16)** You will need to keep track of the number of probes in order to apply quadratic probing.
- R-10.17)** Think of where the entry with minimum key is stored.
- R-10.18)** Since the map will still contain  $n$  entries at the end, you can assume that each `remove()` operation takes the same asymptotic time.
- R-10.19)** Take advantage of the existing `findIndex` method.
- R-10.20)** The crucial methods are `get` and `put`.
- R-10.21)** In the new code, what happens when `key` equals `table.get(mid)`?
- R-10.22)** Assume that a skip list is used to implement the sorted map.
- R-10.23)** Mimic the style of the figures in the book.
- R-10.24)** You must link out the removed entry's tower from all the lists it belongs to.
- R-10.25)** Compare to the implementation of the `addAll` method given in Section 10.5.1.
- R-10.26)** Compare to the implementation of the `addAll` method given in Section 10.5.1.
- R-10.27)** Recall that most skip-list operations run in  $O(\log k)$  time for a set with size  $k$ .
- R-10.28)** Recall that most hash-table operations run in  $O(1)$  expected time.
- R-10.29)** Recall that most hash-table operations run in  $O(1)$  expected time.
- R-10.30)** Recall that most hash-table operations run in  $O(1)$  expected time.
- R-10.31)** Something from this chapter should be helpful!

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## Creativity

- C-10.32)** Consider a bootstrapping method for finding the primes up to  $\sqrt{2M}$ .
- C-10.33)** Your solution should make a single call to `findIndex`.
- C-10.34)** You may assume that such a method is supported by the auxiliary `UnsortedTableMap`
- C-10.35)** You must only call the `findSlot` utility once.
- C-10.36)** Model your solution after the existing support for other map methods.
- C-10.37)** You may assume that such a method is supported by the auxiliary `UnsortedTableMap`

- C-10.38)** The heart of the process can be performed by the findSlot utility.
- C-10.39)** When might the load factor fall below the threshold, and how can you detect this from within AbstractHashMap?
- C-10.40)** Start by defining an appropriate subclass of AbstractMap.MapEntry that includes the new next field that is desired.
- C-10.41)** You need to do some shifting of entries to close up the “gap” just made, but you should only do this for entries that need to move.
- C-10.42)** For part a, note that the symmetry will halve the range of possible values. For part b, note that such automatic collisions will not occur.
- C-10.43)** Perhaps borrow techniques from the Progression hierarchy of Chapter 2, or use Java’s iterators.
- C-10.44)** Maintain a secondary PositionalList instance that represents the FIFO order, and store positions within that list with each entry of the hash table.
- C-10.45)** Each map entry instance should store its current index within the table.
- C-10.46)** Each map entry instance should store its current index within the table.
- C-10.47)** Each map entry instance should store its position within the bucket list.
- C-10.48)** Try out some examples.
- C-10.49)** Do a “double” binary search.
- C-10.50)** Dovetail two binary searches.
- C-10.51)** Do a lazy iteration through indices of the underlying array list.
- C-10.52)** Manage a primary iteration through all nonempty buckets, and then a secondary iteration through each element of a bucket.
- C-10.53)** Do a lazy iteration through the nonempty cells of the table.
- C-10.54)** Think about first sorting the pairs by cost.
- C-10.55)** In the insertion algorithm, first repeatedly flip the coin to determine the level where you should start inserting the new entry.
- C-10.56)** Consider augmenting each node  $v$  in a higher level with the number of missing entries in the gap from  $v$  to the next node over.
- C-10.57)** Consider augmenting each node  $v$  in a higher level with the number of missing entries in the gap from  $v$  to the next node over.
- C-10.58)** Think first about how you can determine the number of 1’s in any row in  $O(\log n)$  time.
- C-10.59)** Consider a two-pass solution, with use of an auxiliary structure.
- C-10.60)** Think of describing your algorithms in terms of boolean operations on the bit vectors.

- C-10.61)** Think of how you could transform  $D$  into  $L$ .
- C-10.62)** Consider the way `subMap` was implemented for `SortedTableMap`.
- C-10.63)** You need some way of grouping together entries with the same key.
- C-10.64)** You need some way of grouping together elements with the same key.

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## Projects

- P-10.65)** The biggest challenge is detecting the case of an infinite loop.
- P-10.66)** When searching for an existing key, make sure to consider both of the possible buckets.
- P-10.67)** Maintain a secondary `PositionalList` instance that represents the FIFO order, and store positions within that list with each entry of the hash table.
- P-10.68)** We've already implemented them for you; all that's left is the experiments.
- P-10.69)** In a Unix/Linux system, a good place to start is `/usr/dict`.
- P-10.70)** It is okay to generate these phone numbers more-or-less at random.
- P-10.71)** Sentinels can be used in place of the theoretical  $-\infty$  and  $+\infty$ .
- P-10.72)** Try to make your screen images mimic the skip list figures in the book.
- P-10.73)** You need to find some way to let the intermediate nodes in the skip list keep track of the number of elements that have been skipped over.
- P-10.74)** For each word  $t$  that results from a minor change to  $s$ , you can test if  $t$  is in  $W$  in  $O(1)$  time.